Pattern Models and Action Models are Incomparable in Update Expressivity

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#### 1. Context Dynamic-Network Models

## Communication is performed in synchronous rounds Adversary

 A set of infinite sequences of communication graphs (Reflexive directed graphs)



#### **Oblivious** adversary

Any communication graph in a given set X may occur in any round

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We say that X is the adversary

1. Context Iterated Immediate Snapshot (IIS)

#### ▶ IIS can be described as an oblivious dynamic-network model

#### Two processes







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#### 1. Context After first round



Dynamic Epistemic Logics (DEL)

Epistemic Logic augmented with update modalities

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Action Model Logic

Epistemic change is defined via events

Dynamic Epistemic Logics (DEL)

- Epistemic Logic augmented with update modalities
- Action Model Logic
  - Epistemic change is defined via events
    - Indistinguishability between events w.r.t. each agent

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What must be true for an event to occur?

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- What must be true for an event to occur?
- Pattern Model Logic
  - Designed for analyzing distributed systems

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    - What must be true for an event to occur?
- Pattern Model Logic
  - Designed for analyzing distributed systems
    - Who communicates with whom? (communication graph)

- full-information communication
- A graph precondition depends on the model

#### 2. Action Models and Pattern Models Languages

Given a set of agents A and a set of propositions P,

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$$\begin{array}{l} \blacktriangleright \ \mathcal{L}_D \\ \bullet \ \phi := p_a \mid \neg \phi \mid \phi \land \phi \mid D_B \phi \\ \bullet \ \mathcal{L}_{\otimes} \\ \bullet \ \phi := p_a \mid \neg \phi \mid \phi \land \phi \mid D_B \phi \mid [\mathsf{U}, \mathsf{e}] \phi \\ \bullet \ \mathcal{L}_{\odot} \\ \bullet \ \phi := p_a \mid \neg \phi \mid \phi \land \phi \mid D_B \phi \mid [\mathcal{P}, G] \phi \end{array}$$

#### 2. Action Models and Pattern Models Action Model

- $\mathsf{U}=(\mathsf{E},\mathsf{R},\mathsf{Pre})$ 
  - E a set of events
  - $R: A \to \wp(\mathsf{E} \times \mathsf{E})$  (indistinguishability)
  - ▶ Pre :  $\mathsf{E} \to \mathcal{L}_D$



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$$M' = (W', \sim', L') = M \otimes \mathsf{U}$$

$$W' = \{(w, \mathbf{e}) \in W \times \mathsf{E} \mid M, w \models \mathsf{Pre}(\mathbf{e})\}$$

$$\sim'_a = \{((w, \mathbf{e}), (w', \mathbf{e}')) \in W' \times W' \mid w \sim_a w' \land \mathbf{e} \mathsf{R}_a \mathbf{e}'\}$$

$$L'((w, \mathbf{e})) = L(w)$$

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#### 2. Action Models and Pattern Models Pattern Model

 $\mathcal{P} = (\mathbf{G}, \mathit{Pre})$ 

 G a set of communication graphs

$$\blacktriangleright Pre: \mathbf{G} \to \mathcal{L}_D$$



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$$W' = \{(w, G) \in W \times \mathbf{G} \mid M, w \models Pre(G) \}$$
  
▶  $\sim'_a = \{((w, G), (w', G')) \in W' \times W' \mid Ga = G'a \land w \sim_{Ga} w'\}$ 

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Ga in-neighborhood of a in G

$$\sim_B = \bigcap_{b \in B} \sim_b$$

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#### 2. Action Models and Pattern Models Semantics on epistemic models

$$\blacktriangleright M, w \models p_a \text{ iff } p_a \in L(w)$$

- $\blacktriangleright M, w \models \neg \phi \text{ iff } M, w \not\models \phi$
- $\blacktriangleright \ M,w\models\phi\wedge\psi \text{ iff }M,w\models\phi \text{ and }M,w\models\psi$

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$$M, w \models D_B \phi$$
 iff  $M, v \models \phi$  for all  $v \sim_B w$ 

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- $\blacktriangleright M, w \models [\mathsf{U},\mathsf{e}]\phi \text{ iff } M, w \models \mathsf{Pre}(\mathsf{e}) \text{ implies } M \otimes \mathsf{U}, (w,\mathsf{e}) \models \phi$

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 $M, w \models [\mathcal{P}, G]\phi \text{ iff } M, w \models Pre(G) \text{ implies } M \odot \mathcal{P}, (w, G) \models \phi$ 

What questions do we want to answer?



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Given an action model U, is there a pattern model P with the same update effect as U?

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- Given a pattern model *P*, is there an action model U with the same update effect as *P*?

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By notational abbreviation,

$$\begin{aligned} \bullet \quad [\mathsf{U}]\phi &:= \bigwedge_{\mathsf{e}\in E} [\mathsf{U},\mathsf{e}]\phi \\ \bullet \quad [\mathcal{P}]\phi &:= \bigwedge_{G\in \mathbf{G}} [\mathcal{P},G]\phi \end{aligned}$$

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$$\phi_0 = 0_a \wedge 0_b$$
  

$$\phi_1 = 0_a \wedge 1_b$$
  

$$\phi_3 = 1_a \wedge 1_b$$
  

$$\phi_2 = 1_a \wedge 0_b$$



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#### Observations





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▶  $M^0 \otimes U$  has 36 worlds

• There are just four communication graphs for  $A = \{a, b\}$ 

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Pattern models are not at least as update expressive as action models

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Given a pattern model  $\mathcal P,$  is there an action model U with the same update effect as  $\mathcal P?$ 

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$$\blacktriangleright M^0(\odot \mathcal{P}_{two-IIS})^n$$



$$(\bullet): \quad (3, G^{ab^{n-1}}G^{b.a}) \xrightarrow{a} (3, G^{ab^n}) \xrightarrow{b} (3, G^{ab^{n-1}}G^{a.b})$$

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Let us assume that there is an action model U with the same update effect as  $\mathcal{P}_{two-IIS}$ 

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Let us assume that there is an action model U with the same update effect as  $\mathcal{P}_{two-I\!I\!S}$ 

The modal depth (md) of U is the maximum modal depth of its precondition formulas

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Close worlds become bounded collective bisimilar

Let us assume that there is an action model U with the same update effect as  $\mathcal{P}_{two-IIS}$ 

The modal depth (md) of U is the maximum modal depth of its precondition formulas

- Close worlds become bounded collective bisimilar
- $M^0(\odot \mathcal{P}_{two-IIS})^{\mathbf{n}+1}$  and  $M^0(\odot \mathcal{P}_{two-IIS})^{\mathbf{n}} \otimes \mathsf{U}$  are not collectively bisimilar.

Let us assume that there is an action model U with the same update effect as  $\mathcal{P}_{two-IIS}$ 

- The modal depth (md) of U is the maximum modal depth of its precondition formulas
- Close worlds become bounded collective bisimilar
- M<sup>0</sup>(⊙P<sub>two-IIS</sub>)<sup>n+1</sup> and M<sup>0</sup>(⊙P<sub>two-IIS</sub>)<sup>n</sup> ⊗ U are not collectively bisimilar.

$$n>\log_3 2(md(\mathsf{U})+1)$$

 $(\cdot, G^{b.a}) \xrightarrow{a} (\cdot, G^{ab}) \xrightarrow{b} (\cdot, G^{a.b}) \xrightarrow{a} (\cdot, G^{a.b}) \xrightarrow{b} (\cdot, G^{a.b}) \xrightarrow{a} (\cdot, G^{b.a}) \xrightarrow{b} (\cdot, G^{b.a}) \xrightarrow{a} (\cdot, G^{a.b}) \xrightarrow{b} (\cdot, G^{a.b}$ 

 $(\cdot, G^{b,a}) \xrightarrow{a} (\cdot, G^{ab}) \xrightarrow{b} (\cdot, G^{a,b}) \xrightarrow{a} (\cdot, G^{a,b}) \xrightarrow{b} (\cdot, G^{ab}) \xrightarrow{a} (\cdot, G^{b,a}) \xrightarrow{b} (\cdot, G^{b,a}) \xrightarrow{a} (\cdot, G^{ab}) \xrightarrow{b} (\cdot, G^{a,b}) \xrightarrow{b} (\cdot, G^{a,b})$ 

 $\begin{array}{c} (G^{a.b}, \theta) & (G^{a.b}, \theta) & (G^{b.a}, \theta) \\ \hline \theta & \theta & \theta \\ \hline \theta & \theta & \theta \\ \end{array}$ 

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 $(\cdot, G^{b,a}) \xrightarrow{a} (\cdot, G^{ab}) \xrightarrow{b} (\cdot, G^{a,b}) \xrightarrow{a} (\cdot, G^{a,b}) \xrightarrow{b} (\cdot, G^{ab}) \xrightarrow{a} (\cdot, G^{b,a}) \xrightarrow{b} (\cdot, G^{b,a}) \xrightarrow{a} (\cdot, G^{ab}) \xrightarrow{b} (\cdot, G^{a,b}) \xrightarrow{b} (\cdot, G^{a,b})$ 

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 $(\cdot, G^{b,a}) \xrightarrow{a} (\cdot, G^{ab}) \xrightarrow{b} (\cdot, G^{a,b}) \xrightarrow{a} (\cdot, G^{a,b}) \xrightarrow{b} (\cdot, G^{ab}) \xrightarrow{a} (\cdot, G^{b,a}) \xrightarrow{b} (\cdot, G^{b,a}) \xrightarrow{a} (\cdot, G^{ab}) \xrightarrow{b} (\cdot, G^{a,b}) \xrightarrow{b} (\cdot, G^{a,b})$ 

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There is a shorter path to the worlds with different labeling above

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 $(\cdot, G^{b,a}) \xrightarrow{a} (\cdot, G^{ab}) \xrightarrow{b} (\cdot, G^{a,b}) \xrightarrow{a} (\cdot, G^{a,b}) \xrightarrow{b} (\cdot, G^{ab}) \xrightarrow{a} (\cdot, G^{b,a}) \xrightarrow{b} (\cdot, G^{b,a}) \xrightarrow{a} (\cdot, G^{ab}) \xrightarrow{b} (\cdot, G^{a,b}) \xrightarrow{b} (\cdot, G^{a,b})$ 

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There is a shorter path to the worlds with different labeling above

## Action models are not at least as update expressive as pattern models

## Action models and pattern models are incomparable in update expressivity



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